

Quantum Computing

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HW 5: Model Inference

- This homework does **NOT** have sample code
- Write a program for handwritten digit dataset
- Output a file with predictions for each data instance
- You are free to use any model or data to train your model offline
- We will only do model inference during the test
- Two scripts are mandatory:
 - `compile.sh`
 - `run.sh <test_dataset> <output_file>`

HW 5: Model Inference

- The testing dataset is a variation of mnist.t (added gaussian noise)
- We will have 10 cases where we vary the level of noises
 - $N(0,0)$, $N(0,1)$, $N(0,2)$, ... $N(0,9)$
 - $N(0,0)$ corresponds to no noises, i.e., plain mnist.t
- Each test case contains 10k instances.
- A test case is considered correct if the test accuracy is no less than **50%**

HW 5: Model Inference

- No 3rd party code is allowed.
- 10 test cases. Each case weights 1 pt.
- The compilation is considered failed if it does not finish in **5 minute**.
- A test case is considered **incorrect** if it does not finish in **2 minutes**.
- **Correct GPU solutions will get 5 pts bonus.**
- The **summation** of the execution time across 10 cases will be used to rank **correct** solutions.

- Due: 11/30/2022 **1:00 pm** EST

Testing Environment

- `ssh yourusername@granger.cs.rit.edu`
- Intel(R) Xeon(R) CPU E5-2650 v4 @ 2.20GHz
- 48 threads in total (2 sockets, 12 cores per socket, 2 threads per core)
- 251 GB memory
- GPU: Tesla P4
- Testing limit:
 - 8 threads `taskset -c`
 - 2 GPU

Quantum Gates

- Not
- Hadamard
- Controlled Not

- No Cloning
- No Forgetting

- Entangling

Deutsch–Jozsa algorithm

$$\frac{1}{2}(|0\rangle + |1\rangle)(|0\rangle - |1\rangle) \quad f \text{ that maps } |x\rangle|y\rangle \text{ to } |x\rangle|f(x) \oplus y\rangle$$

$$\frac{1}{2}(|0\rangle(|f(0) \oplus 0\rangle - |f(0) \oplus 1\rangle) + |1\rangle(|f(1) \oplus 0\rangle - |f(1) \oplus 1\rangle))$$

$$= \frac{1}{2}((-1)^{f(0)}|0\rangle(|0\rangle - |1\rangle) + (-1)^{f(1)}|1\rangle(|0\rangle - |1\rangle))$$

$$= (-1)^{f(0)} \frac{1}{2} \left(|0\rangle + (-1)^{f(0) \oplus f(1)} |1\rangle \right) (|0\rangle - |1\rangle).$$

$$\begin{aligned} \frac{1}{\sqrt{2}}(|0\rangle + (-1)^{f(0) \oplus f(1)} |1\rangle) & \quad \frac{1}{2}(|0\rangle + |1\rangle + (-1)^{f(0) \oplus f(1)} |0\rangle - (-1)^{f(0) \oplus f(1)} |1\rangle) \\ & = \frac{1}{2}((1 + (-1)^{f(0) \oplus f(1)})|0\rangle + (1 - (-1)^{f(0) \oplus f(1)})|1\rangle) \end{aligned}$$

CHSH Game

- John Clauser, Michael Horne, Abner Shimony, and Richard Holt in 1969
- C: Random generate x, y
- A gets x and generates a
- B gets y and generates b
- A and B wins if $x * y = a + b$

Shor's Algorithm

1. Pick a random number $1 < a < N$.
2. Compute $K = \gcd(a, N)$, the [greatest common divisor](#) of a and N . This may be done using the [Euclidean algorithm](#).
3. If $K \neq 1$, then K is a [nontrivial](#) factor of N , so we are done.
4. Otherwise, use the quantum period-finding subroutine (below) to find r , which denotes the [period](#) of the following function:

$$f(x) = a^x \bmod N.$$

Equivalently, r is the smallest positive integer that satisfies $a^r \equiv 1 \pmod N$.

5. If r is odd, then go back to step 1.
6. If $a^{r/2} \equiv -1 \pmod N$, then go back to step 1.
7. Otherwise, both $\gcd(a^{r/2} + 1, N)$ or $\gcd(a^{r/2} - 1, N)$ are nontrivial factors of N , so we are done.