



CSCI 740 - Programming Language Theory

Lecture 30

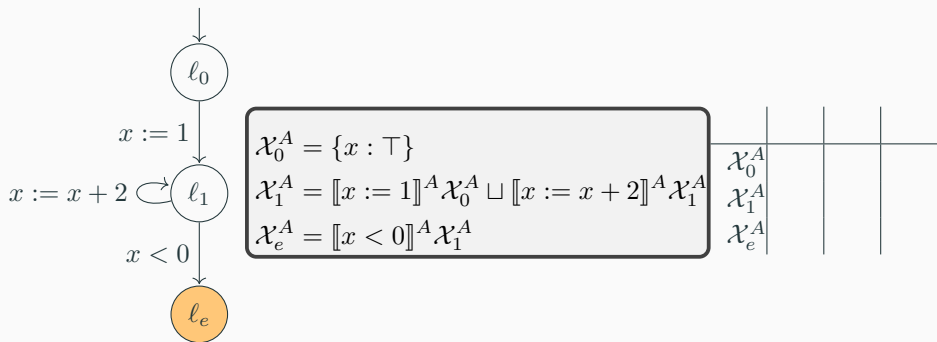
Widening and Narrowing

Instructor: Hossein Hojjat

November 13, 2017

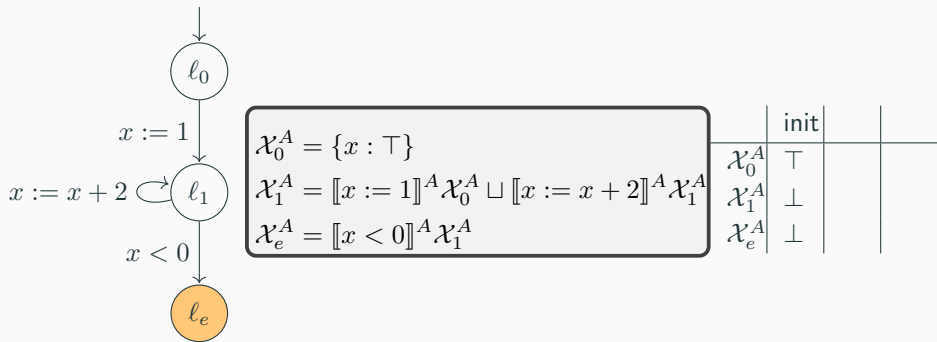
Example: Sign Abstraction

Use sign abstraction to prove error is unreachable $\text{Abs} = \{\top, +, 0, -, \perp\}$



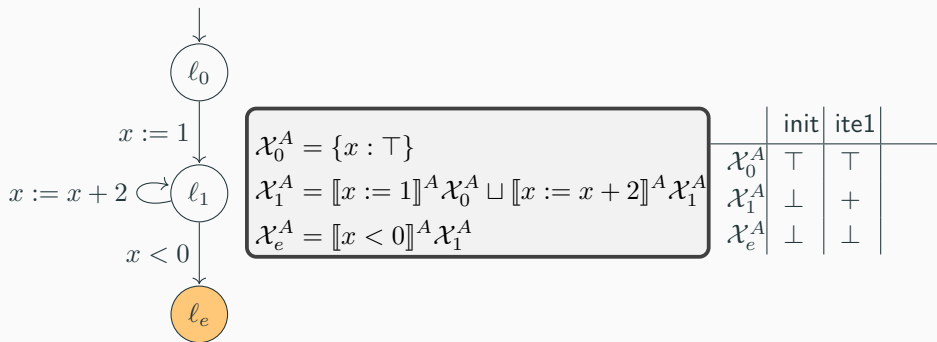
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Example: Sign Abstraction

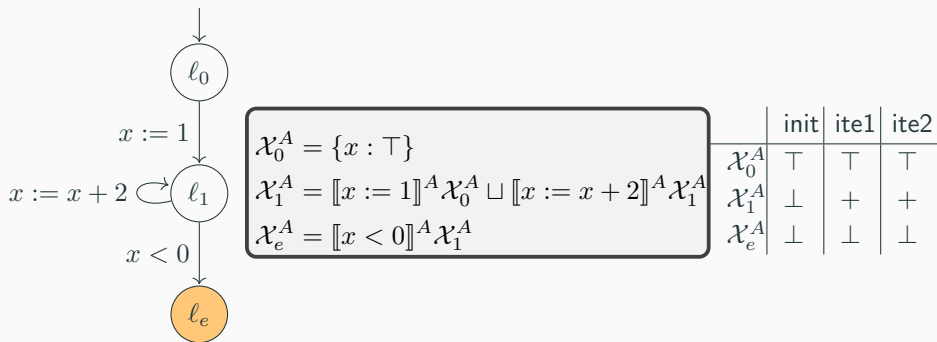
Use sign abstraction to prove error is unreachable $\text{Abs} = \{\top, +, 0, -, \perp\}$



$$\begin{aligned}\mathcal{X}_1^A &= \llbracket x := 1 \rrbracket^A \mathcal{X}_0^A \sqcup \llbracket x := x + 2 \rrbracket^A \mathcal{X}_1^A \\ &= \llbracket x := 1 \rrbracket^A \{x : \top\} \sqcup \llbracket x := x + 2 \rrbracket^A \{x : \perp\} \\ &= \{x : +\} \sqcup \{x : \perp\} \\ &= \{x : + \sqcup \perp\} \\ &= \{x : +\}\end{aligned}$$

Example: Sign Abstraction

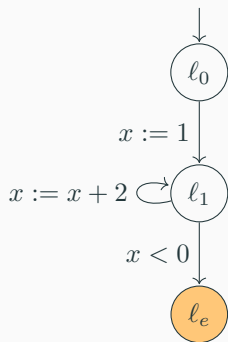
Use sign abstraction to prove error is unreachable $\text{Abs} = \{\top, +, 0, -, \perp\}$



$$\begin{aligned}\mathcal{X}_1^A &= \llbracket x := 1 \rrbracket^A \mathcal{X}_0^A \sqcup \llbracket x := x + 2 \rrbracket^A \mathcal{X}_1^A \\ &= \llbracket x := 1 \rrbracket^A \{x : \top\} \sqcup \llbracket x := x + 2 \rrbracket^A \{x : +\} \\ &= \{x : +\} \sqcup \{x : +\} \\ &= \{x : + \sqcup +\} \\ &= \{x : +\}\end{aligned}$$

Example: Sign Abstraction

Use sign abstraction to prove error is unreachable $\text{Abs} = \{\top, +, 0, -, \perp\}$



$$\mathcal{X}_0^A = \{x : \top\}$$

$$\mathcal{X}_1^A = \llbracket x := 1 \rrbracket^A \mathcal{X}_0^A \sqcup \llbracket x := x + 2 \rrbracket^A \mathcal{X}_1^A$$

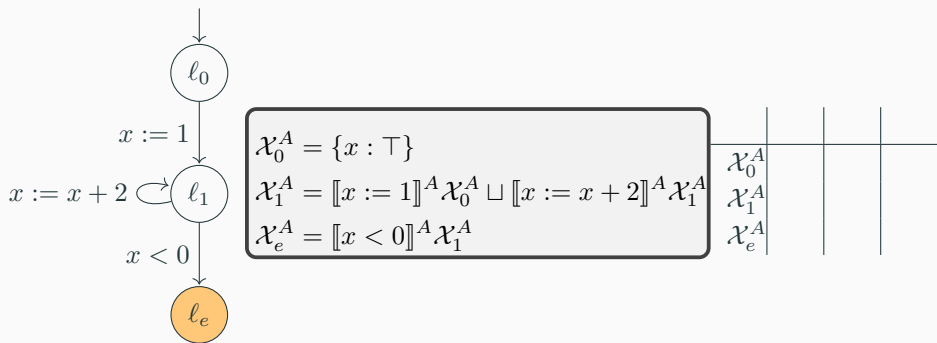
$$\mathcal{X}_e^A = \llbracket x < 0 \rrbracket^A \mathcal{X}_1^A$$

	init	ite1	ite2
\mathcal{X}_0^A	\top	\top	\top
\mathcal{X}_1^A	\perp	$+$	$+$
\mathcal{X}_e^A	\perp	\perp	\perp

fixpoint - error unreachable

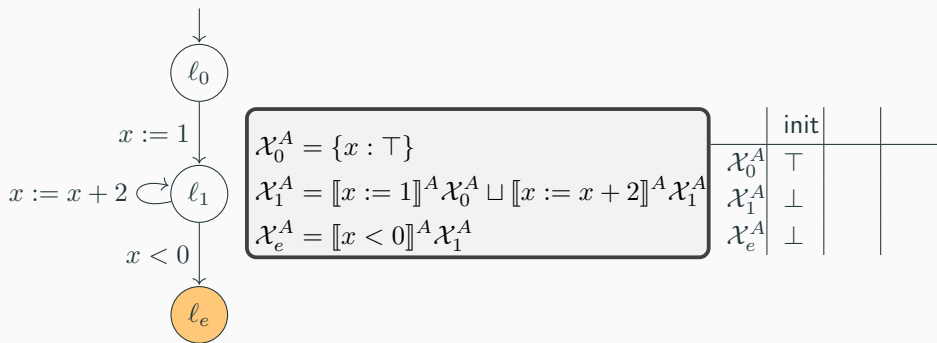
Exercise: Interval Abstraction

Use interval abstraction to prove error is unreachable



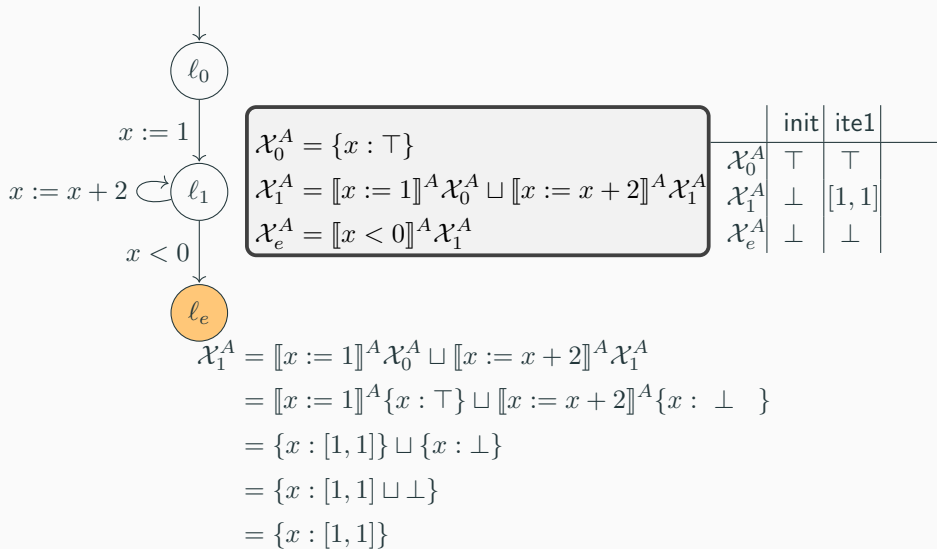
Exercise: Interval Abstraction

Use interval abstraction to prove error is unreachable



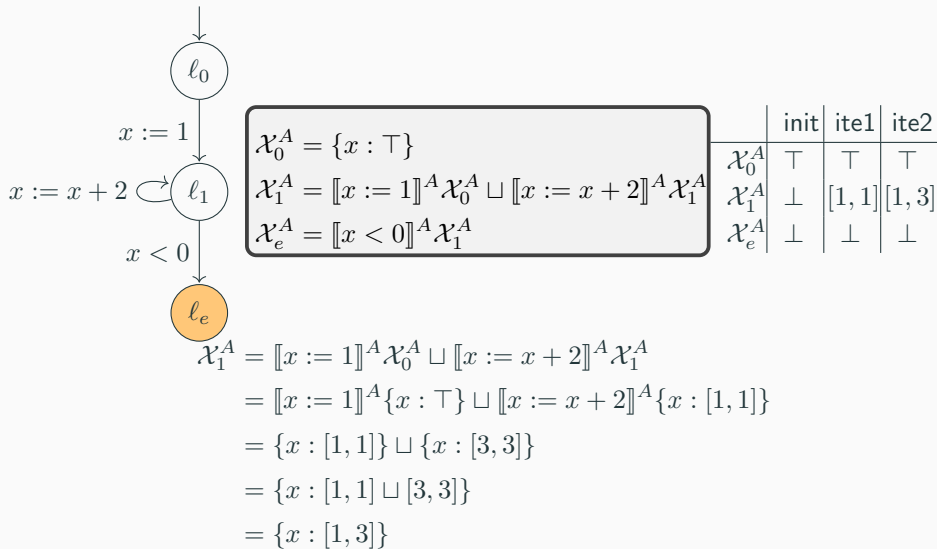
Exercise: Interval Abstraction

Use interval abstraction to prove error is unreachable



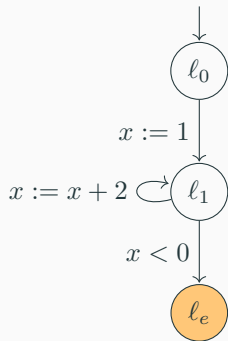
Exercise: Interval Abstraction

Use interval abstraction to prove error is unreachable



Exercise: Interval Abstraction

Use interval abstraction to prove error is unreachable



$$\mathcal{X}_0^A = \{x : \top\}$$

$$\mathcal{X}_1^A = \llbracket x := 1 \rrbracket^A \mathcal{X}_0^A \sqcup \llbracket x := x + 2 \rrbracket^A \mathcal{X}_1^A$$

$$\mathcal{X}_e^A = \llbracket x < 0 \rrbracket^A \mathcal{X}_1^A$$

	init	ite1	ite2
\mathcal{X}_0^A	\top	\top	\top
\mathcal{X}_1^A	\perp	$[1, 1]$	$[1, 3]$
\mathcal{X}_e^A	\perp	\perp	\perp

... will go on forever

$$\mathcal{X}_1^A : \perp, [1, 1], [1, 3], [1, 5], \dots$$

Widening & Narrowing

- Some interesting abstract domains have infinite lattices
 - Interval domain
- These domains have infinite ascending chains:
 - $[0, 0], [0, 1], [0, 2], \dots$
- Abstraction interpretation usually provides only simpler computation but not convergence
- In some cases we may need to accelerate the fixpoint computation using widening
- If we do too much widening then we may need narrowing

Widening

- Intuition:
widening operator $\nabla : \mathcal{D} \rightarrow \mathcal{D}$ jumps ahead on a lattice $\langle \mathcal{D}, \leq \rangle$
- ∇ is like join \sqcup but can be a bigger element of the lattice

$$x \leq x \nabla y$$

$$y \leq x \nabla y$$

- For all increasing chains $x_0 \leq x_1 \leq \dots$, the increasing chain (y_i) defined as

$$y_i = \begin{cases} x_0 & \text{if } i = 0; \\ y_{i-1} \nabla x_i & \text{if } i > 0. \end{cases}$$

eventually stabilizes (i.e., the chain is finite)

- In other words, iterated uses of widening operator becomes a fixpoint eventually

$$x_0 \nabla x_1 \nabla x_2 \nabla x_3 \dots$$

- Jump ahead to infinity whenever interval bounds are not inclusive

$$[a, b] \nabla \perp = [a, b]$$

$$\perp \nabla [c, d] = [c, d]$$

$$[a, b] \nabla [c, d] = \left[\left\{ \begin{array}{ll} a & \text{if } a \leq c \\ -\infty & \text{otherwise} \end{array} \right\}, \left\{ \begin{array}{ll} b & \text{if } d \leq b \\ +\infty & \text{otherwise} \end{array} \right\} \right]$$

Non-monotonicity

- Non monotonicity of widening
- $[0, 1] \nabla [0, 2] =$
- $[0, 2] \nabla [0, 2] =$

- Jump ahead to infinity whenever interval bounds are not inclusive

$$[a, b] \nabla \perp = [a, b]$$

$$\perp \nabla [c, d] = [c, d]$$

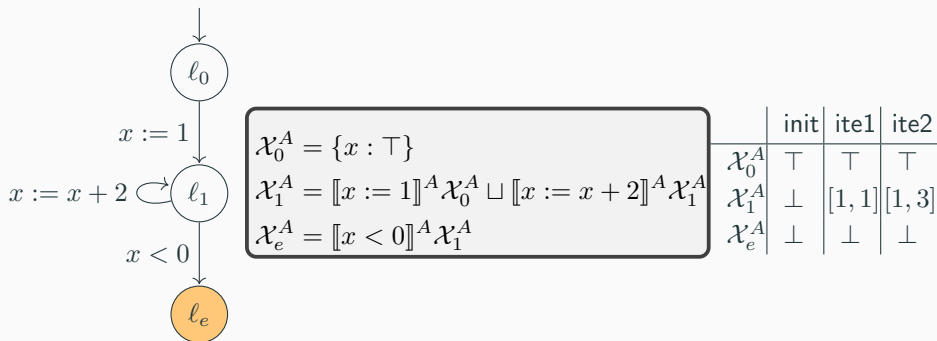
$$[a, b] \nabla [c, d] = \left[\left\{ \begin{array}{ll} a & \text{if } a \leq c \\ -\infty & \text{otherwise} \end{array} \right\}, \left\{ \begin{array}{ll} b & \text{if } d \leq b \\ +\infty & \text{otherwise} \end{array} \right\} \right]$$

Non-monotonicity

- Non monotonicity of widening
- $[0, 1] \nabla [0, 2] = [0, +\infty)$
- $[0, 2] \nabla [0, 2] = [0, 2]$

Interval Abstraction

Use interval abstraction (with widening) to prove error is unreachable



$$\mathcal{X}_0^A = \top$$

$$\mathcal{X}_1^A = [1, 1] \nabla [1, 3] = [1, +\infty)$$

$$\mathcal{X}_e^A = \perp$$

Iterations stabilize to a fix-point

Narrowing

- When widening is too aggressive we use a dual operator called narrowing
- Δ is like meet \sqcap but can be a bigger element of the lattice

$$x \sqcap y \leq x \Delta y$$

- For all increasing chains $x_0 \leq x_1 \leq \dots$, the increasing chain (y_i) defined as

$$y_i = \begin{cases} x_0 & \text{if } i = 0; \\ y_{i-1} \Delta x_i & \text{if } i > 0. \end{cases}$$

eventually stabilizes (i.e., the chain is finite)

Narrowing for Intervals

$$[a, b] \Delta \perp = \perp$$

$$\perp \Delta [c, d] = \perp$$

$$[a, b] \Delta [c, d] = \left[\left\{ \begin{array}{ll} c & \text{if } a = -\infty \\ a & \text{otherwise} \end{array} \right\}, \left\{ \begin{array}{ll} d & \text{if } b = +\infty \\ b & \text{otherwise} \end{array} \right\} \right]$$