



CSCI-344

Programming Language Concepts (Section 3)

Lecture 2

Impcore: an Imperative Core Language

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Programming Paradigm

- A **programming paradigm** is a model of programming based on concepts that shapes the way programmers think, design, organize and write programs
- Main programming paradigms:
 - Imperative programming
 - Functional programming
 - Logic programming
- Orthogonal to it:
 - Object-oriented programming

Programming Paradigm

- A **programming paradigm** is a model of programming based on concepts that shapes the way programmers think, design, organize and write programs
- Main programming paradigms:
 - Imperative programming ← Today's lecture
 - Functional programming
 - Logic programming
- Orthogonal to it:
 - Object-oriented programming

- **Imperative programming**

- Modifying mutable variables using assignments
- Sequence of statements, with side effects
- Imperative: “do this, then do that, then do this other thing, ...”
- Control structures if-branches, loops, break, continue, etc

- **Imperative programming**

- Modifying mutable variables using assignments
- Sequence of statements, with side effects
- Imperative: “do this, then do that, then do this other thing, ...”
- Control structures if-branches, loops, break, continue, etc

- **ImpCore (IMPerative CORE)**

- Simple imperative programming language
- Embodies a tiny core fragment of conventional mainstream languages such as C++, Java, ...
- Eliminate superficial differences, focus on main features

Example to find number is prime

```
(val prime 1)
(val i 2)
(define is-prime? (n)
  (begin
    (while (<= i (/ n 2))
      (begin
        (if (= (mod n i) 0)
          (set prime 0)
          (set prime prime) ;; don't change
        )
        (set i (+ i 1))
      )
    )
    prime
  )
)
(print (is-prime? 31))
```

Example to find number is prime

```
(val prime 1)
```

```
(val i 2)
```

```
(define is-prime? (n)
```

```
  (begin
```

```
    (while (<= i (/ n 2))
```

```
      (begin
```

```
        (if (= (mod n i) 0)
```

```
          (set prime 0)
```

```
          (set prime prime) ;; don't change
```

```
        )
```

```
        (set i (+ i 1))
```

```
      )
```

```
    )
```

```
    prime
```

```
  )
```

```
)
```

```
(print (is-prime? 31))
```

Global Variable Definition

Example to find number is prime

```
(val prime 1)
```

```
(val i 2)
```

Function Definition

```
(define is-prime? (n)
  (begin
    (while (<= i (/ n 2))
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        (if (= (mod n i) 0)
          (set prime 0)
          (set prime prime) ;; don't change
        )
        (set i (+ i 1))
      )
    )
    prime
  )
)
```

```
(print (is-prime? 31))
```


Example to find number is prime

```
(val prime 1)
(val i 2)
(define is-prime? (n)                                     Begin Expression
  (begin
    (while (<= i (/ n 2))
      (begin
        (if (= (mod n i) 0)
          (set prime 0)
          (set prime prime) ;; don't change
        )
        (set i (+ i 1))
      )
    )
    prime ←————— Return Value
  )
)
(print (is-prime? 31))
```

Example to find number is prime

```
(val prime 1)
(val i 2)
(define is-prime? (n)
  (begin
    (while (<= i (/ n 2))
      (begin
        (if (= (mod n i) 0)
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          (set prime prime) ;; don't change
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        (set i (+ i 1))
      )
    )
    prime
  )
)
(print (is-prime? 31))
```

While Expression (returns 0)

Example to find number is prime

```
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(define is-prime? (n)
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    (while (<= i (/ n 2))
      (begin
        (if (= (mod n i) 0)
          (set prime 0)
          (set prime prime) ;; don't change
        )
        (set i (+ i 1))
      )
    )
    prime
  )
)
(print (is-prime? 31))
```

If Expression

Prefix notation: $(+ (* x y) 1)$

Infix notation: $(x * y) + 1$

- ImpCore has fully parenthesized prefix notation (as in Lisp)
- Variables and names are basic atoms
- Other constructs bracketed with parentheses
(Possible keyword after opening bracket)
- **Pros:** Trivial to parse
- **Cons:** Unreadable in complex expressions

- Two most important syntactic categories:
definitions and **expressions**
- ImpCore has no statements
- An ImpCore program is a sequence of definitions

ImpCore Definitions

- `(val x exp)`
- `(define f (x ... y) exp)`
- `exp`

Global variable definition (no local variables)

- Example: `(val prime 1)`
- Equivalent construct in C: `int prime = 1;`

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Function definition

- Example: `(define max (x y) (if (> x y) x y))`
- Equivalent construct in C:

```
int max(int x, int y) {  
    return (x>y) ? x : y;  
}
```

Global variable definition (no local variables)

- Example: `(val prime 1)`
- Equivalent construct in C: `int prime = 1;`

Function definition

- Example: `(define max (x y) (if (> x y) x y))`
- Equivalent construct in C:

```
int max(int x, int y) {  
    return (x>y) ? x : y;  
}
```

Top-level expression definition

- Top-level expression definitions are treated as implicit definition of a standard variable `it`
- Example: `(+ 5 6)`
- Equivalent definition: `(val it (+ 5 6))`

ImpCore Expressions

- Expression produces a value and may have side effects
- 100 (integer literal)
- x (variable)
- (**if** exp1 exp2 exp3)
- (**while** exp1 exp2)
- (**set** x exp)
- (**begin** exp1 ... expn)
- (f exp1 ... expn)
- Primitive Functions: + , - , * , / , = , < , > , print
- Only type of data is “machine integer” (deliberate oversimplification)

Concrete vs. abstract syntax

Abstract Syntax

- The “essential” constructs of a programming language

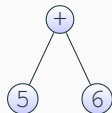
Concrete Syntax

- The actual syntax of a programming language

Concrete vs. abstract syntax

Abstract Syntax

- The “essential” constructs of a programming language
- An additive expression has two operand expressions as its “essential” parts



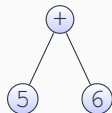
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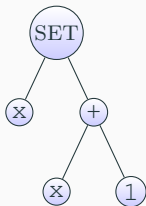
Concrete Syntax

- The actual syntax of a programming language
- | | |
|--------------------------------|------------------|
| - 5 + 6 | Java (infix) |
| - (+ 5 6) | ImpCore (prefix) |
| - bipush 5
bipush 4
iadd | JVM (postfix) |

Concrete vs. abstract syntax

Abstract Syntax

- The “essential” constructs of a programming language



Concrete Syntax

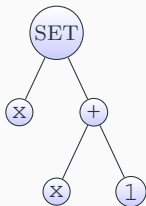
- The actual syntax of a programming language

- <code>x = x + 1;</code>	Java
- <code>(set x (+ x 1))</code>	ImpCore
- <code>x := x + 1;</code>	Pascal

Concrete vs. abstract syntax

Abstract Syntax

- The “essential” constructs of a programming language



Concrete Syntax

- The actual syntax of a programming language

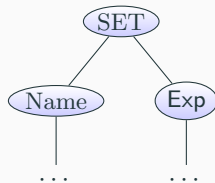
- `x = x + 1;` Java
- `(set x (+ x 1))` ImpCore
- `x := x + 1;` Pascal

Abstract Syntax Tree (AST): Ordered-tree representation of the abstract syntactic structure of source code

Abstract Syntax Tree (AST) Representation

- We represent AST with context-free grammars
- Label each node with all capitalized constructor
- Specify children in parentheses

Exp = LITERAL (Value)
| VAR (Name)
| SET (Name, Exp)
| ...

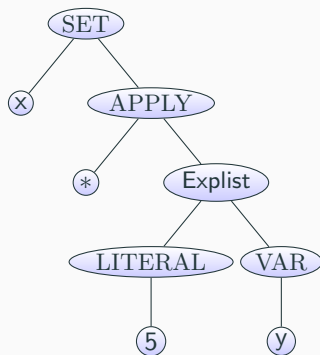


ImpCore Abstract Syntax

Def = VAL (Name, Exp)
| EXP (Exp)
| DEFINE (Name, Namelist, Exp)

Exp = LITERAL (Value)
| VAR (Name)
| SET (Name, Exp)
| IF (Exp, Exp, Exp)
| WHILE (Exp, Exp)
| BEGIN (Explist)
| APPLY (Name, Explist)

(set x (* 5 y))



Typical Interpreter

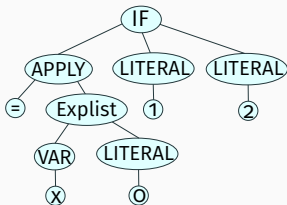
Source Code
(concrete syntax)

```
( if ( = x 0 ) 1 2 )
```

Token Stream

```
( if ( = x 0 ) 1 2 )
```

Abstract Syntax Tree



Lexical Analysis

Syntax Analysis
(Parsing)

Semantic Analysis

Error

- Syntax analysis does not correlate variable definitions with the uses of variables
- Example: Syntactically correct program

```
(val x 0)
```

```
(set x (+ y 1))
```

Semantic Errors

- Syntax analysis does not correlate variable definitions with the uses of variables
- Example: Syntactically correct program

```
(val x 0)
(set x (+ y 1))
```
- `(+ y 1)` cannot be evaluated by itself: we need to know what `y` is
- Programming-language theory uses **environments** to define the meanings of names
- Environment is a mapping from names to meanings
- Environments are often implemented as hash tables or search trees, so they are sometimes called **symbol tables**

- An environment is a set of bindings denoted \mapsto
- We write $\rho\{x \mapsto v\}$ to mean the environment ρ extended by a binding of the name x to v

$$\rho\{x \mapsto v\}(y) = \begin{cases} v, & \text{when } y = x \\ \rho(y), & \text{when } y \neq x \end{cases}$$

- ξ : global variables, ρ : formal parameters
- There are no local variables: if you need temporaries use formal parameters

Example to find number is prime

$\xi_0 = \{\}$

```
(val prime 1)
(val i 2)
(define is-prime? (n)
  (begin
    (while (<= i (/ n 2))
      (begin
        (if (= (mod n i) 0)
          (set prime 0)
          (set prime prime))
        )
      (set i (+ i 1))
    )
  )
  prime
)
(print (is-prime? 31))
```

Example to find number is prime

$\xi_0 = \{\}$
 $\xi_1 = \xi_0\{prime \mapsto 1\}$

```
→ (val prime 1)
   (val i 2)
   (define is-prime? (n)
     (begin
       (while (<= i (/ n 2))
         (begin
           (if (= (mod n i) 0)
               (set prime 0)
               (set prime prime))
           )
         (set i (+ i 1))
       )
     )
     prime
  )
  )
  (print (is-prime? 31))
```

Example to find number is prime

```

                                 $\xi_0 = \{\}$ 
  (val prime 1)                  $\xi_1 = \xi_0\{prime \mapsto 1\}$ 
→ (val i 2)                      $\xi_2 = \xi_1\{i \mapsto 2\}$ 
  (define is-prime? (n)
    (begin
      (while (<= i (/ n 2))
        (begin
          (if (= (mod n i) 0)
            (set prime 0)
            (set prime prime))
          )
          (set i (+ i 1))
        )
      )
    )
    prime
  )
)
(print (is-prime? 31))
```

Example to find number is prime

```

                                 $\xi_0 = \{\}$ 
(val prime 1)                    $\xi_1 = \xi_0\{prime \mapsto 1\}$ 
(val i 2)                        $\xi_2 = \xi_1\{i \mapsto 2\}$ 
→ (define is-prime? (n)          $\xi_2, \rho_0 = \{n \mapsto 31\}$ 
    (begin
      (while (<= i (/ n 2))
        (begin
          (if (= (mod n i) 0)
              (set prime 0)
              (set prime prime))
          )
          (set i (+ i 1))
        )
      )
    )
  )
prime
)
(print (is-prime? 31))
```


Environmental Abuse

- The environment ϕ contains the definition of functions
- Abuse of separate environments

```
-> (val x 1)
```

```
1
```

```
-> (define x (y) (- x y))
```

```
-> (define y (x) (x x))
```

```
-> (y 1)
```

```
0
```

- Which `xs` correspond to globals?
- Which `xs` correspond to formals?
- Which `xs` correspond to functions?

- We study **Semantics**: the meaning of a program, or how it behaves
- What is the meaning of $(+ \ x \ y)$?
- Semantics helps to understand a particular language
 - what you can depend on as a programmer
 - what you must provide as a compiler writer