



# CSCI 742 - Compiler Construction

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## Lecture 7

Regular Expression - Automata Conversion

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February 6, 2017

# Lexer Automatic Construction: Big Picture

## **Input: Token Spec**

- List of regular expressions (RE) in priority order

## **Output: Lexer**

- Reads an input stream and breaks it up into tokens according to REs

## **Algorithm**

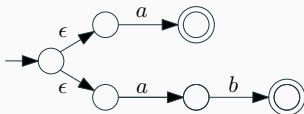
- Convert REs into non-deterministic finite automata (NFA)
- Convert NFA to DFA
- Convert DFA into transition table

# Lexer Automatic Construction: Example

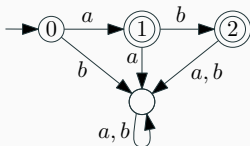
- RE for tokens:

$(a|ab)$

- NFA:



- DFA:



- Transition Table:

	a	b
0	1	Error
1	Error	2
2	Error	Error

## Theorem

A language  $L$  can be described by regular expression if and only if  $L$  is the language accepted by a finite automaton.

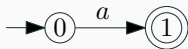
Algorithms:

- regular expression  $\Rightarrow$  automaton    important for lexer construction
- automaton  $\Rightarrow$  regular expression    interesting method in formal languages theory

# RE $\Rightarrow$ Finite Automaton

- Build the finite automaton inductively, based on the definition of regular expressions

$a$



$\epsilon$

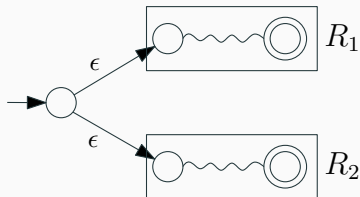


$\emptyset$

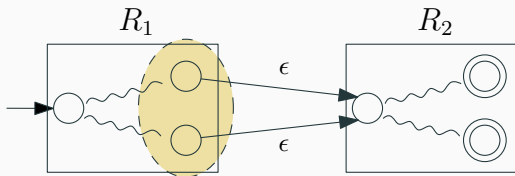


# RE $\Rightarrow$ Finite Automaton

Alternation  $R_1 \mid R_2$

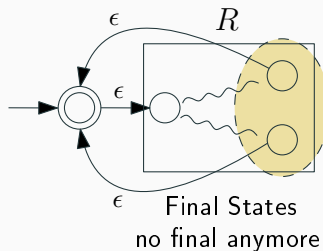


Concatenation  
 $R_1 \cdot R_2$



Final States  
no final anymore

Alternation  $R^*$



## Question

- Construct an NFA for the regular expression  $(ab)^* \mid b^*$

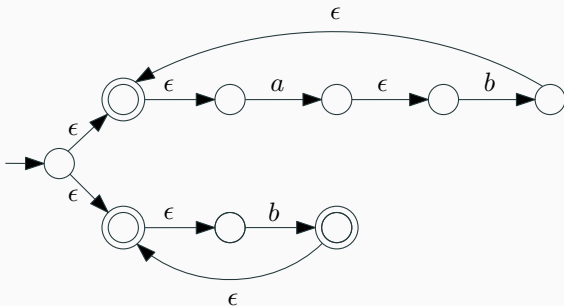


# Exercise

## Question

- Construct an NFA for the regular expression  $(ab)^* \mid b^*$

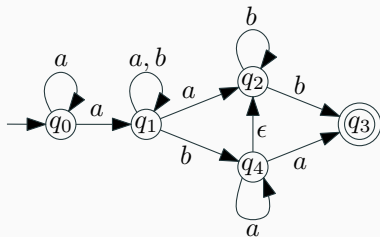
## Answer



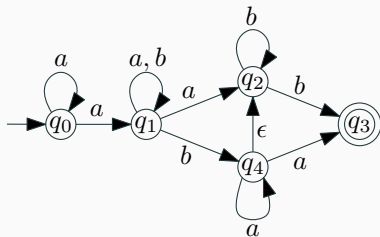
# From NFA to DFA

- For every NFA there exists an equivalent DFA that accepts the same set of strings
- NFAs could be exponentially smaller (succinct)
- Idea: keep track of a set of all possible states in which the automaton could be
- View this finite set as one state of new automaton

# From NFA to DFA: Example

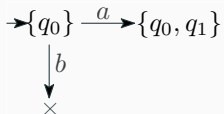
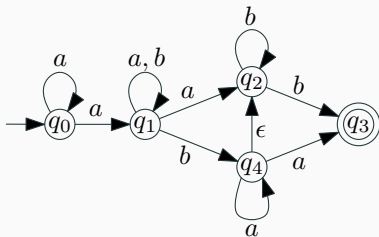


# From NFA to DFA: Example

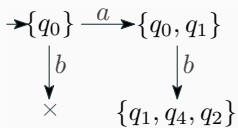
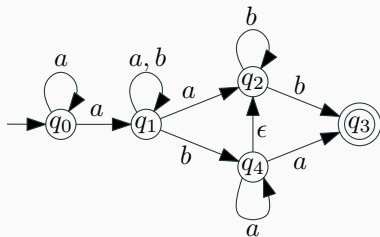


$\rightarrow\{q_0\}$

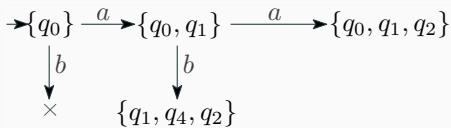
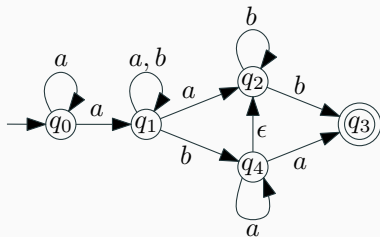
# From NFA to DFA: Example



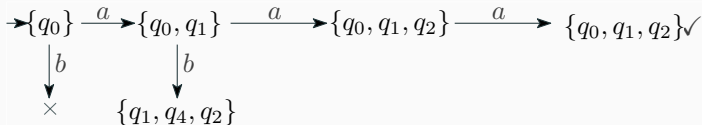
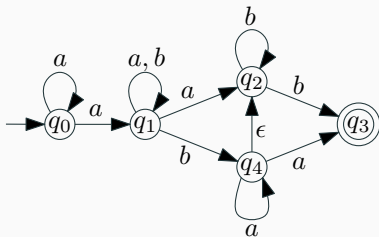
# From NFA to DFA: Example



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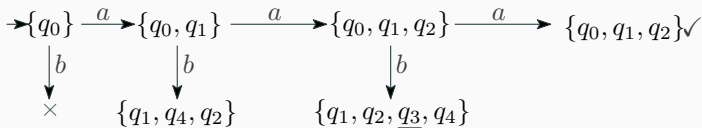
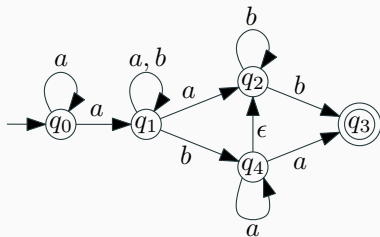
## From NFA to DFA: Example



- When processing if we see a set exactly the same as a set constructed earlier we mark it

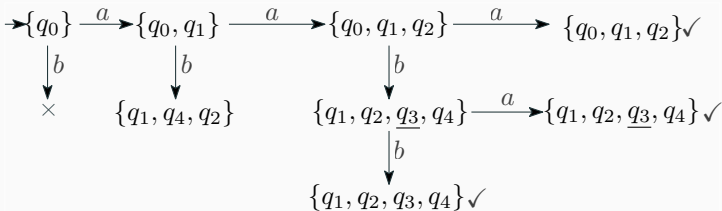
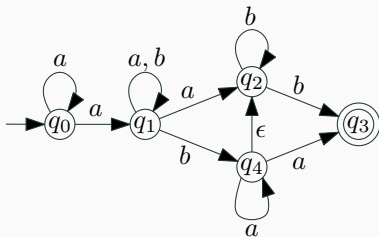


## From NFA to DFA: Example



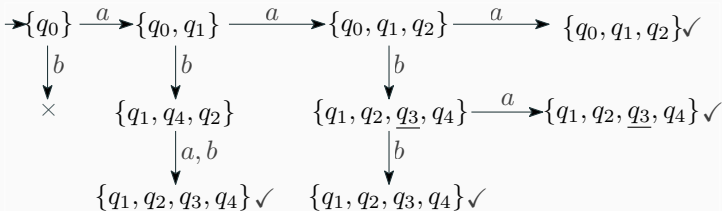
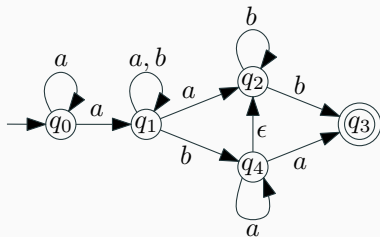
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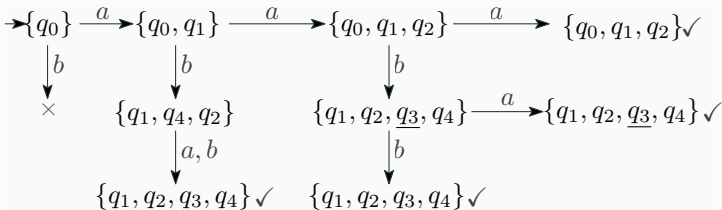
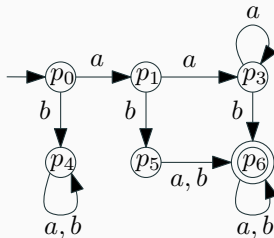
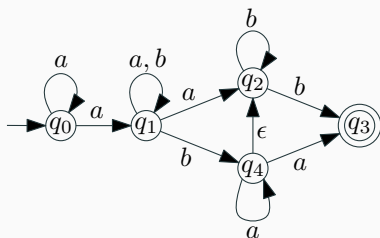
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# From NFA to DFA: Example



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# From NFA to DFA: Example



- When processing if we see a set exactly the same as a set constructed earlier we mark it

## Question

- Construct an NFA for the regular expression  $(bb^*)|a^*$  over alphabet  $\{a, b\}$  and determinize it.