CSCI 742 - Compiler Construction

Lecture 18

LR(0) Parsing

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Recap: Action Selection Problem

- Given stack σ and look-ahead symbol b, should parser:
- **shift** b onto the stack (making it σb)
- **reduce** some production $X \to \gamma$ assuming stack has the form $\alpha\gamma$ (making it αX)

LR Parsing Engine

- Use a set of parser states
- Use stack with alternating symbols and states
  - For example: (3 a 10 + 5 (red = state numbers)
- Use parsing table to:
  - Determine what action to apply (shift/reduce)
  - Determine next state
- The parser actions can be precisely determined from the table

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	reduce 4
0 (1 T 4	) $	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	reduce 4
0 (1 T 4	) $	reduce 2
0 (1 E 6	) $	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	reduce 4
0 (1 T 4	) $	reduce 2
0 (1 E 6	) $	shift 8
0 (1 E 6) 8	$	
		...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$

Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	reduce 4
0 (1 T 4	) $	reduce 2
0 (1 E 6	) $	shift 8
0 (1 E 6) 8	$	reduce 5
0 T 4	$	...

LR(1) Parsing Table Example

	(	)	+	num	$	E	S	T
0	$s1$			$s5$		$g2$	$g3$	$g4$
1	$s1$			$s5$		$g6$		$g4$
2			$s7$		$r1$			
3					acc			
4		$r2$	$r2$		$r2$			
5		$r4$	$r4$		$r4$			
6		$s8$	$s7$					
7	$s1$			$s5$				$g9$
8		$r5$	$r5$		$r5$			
9		$r3$	$r3$		$r3$			

r1	$S \to E\$$
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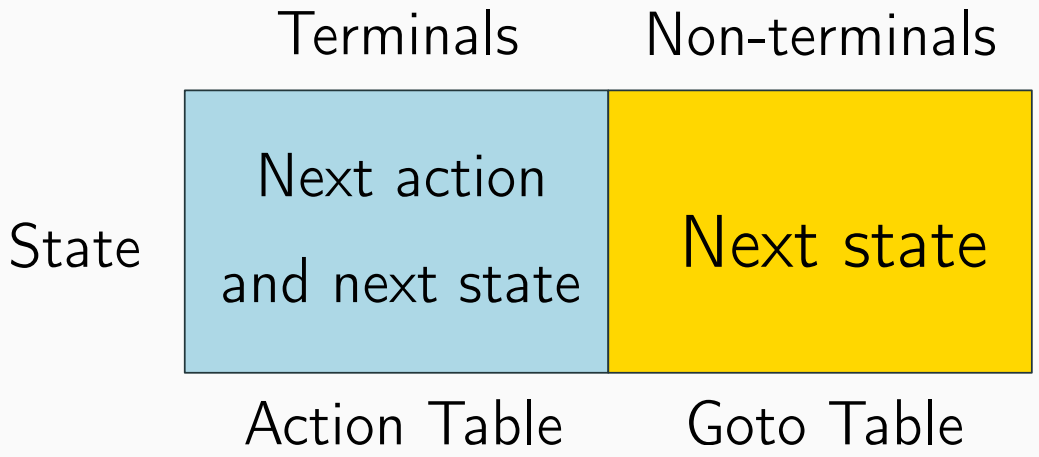
Stack	Input	Action
0	(num) $	shift 1
0 (1	num) $	shift 5
0 (1 num 5	) $	reduce 4
0 (1 T 4	) $	reduce 2
0 (1 E 6	) $	shift 8
0 (1 E 6) 8	$	reduce 5
0 T 4	$	...

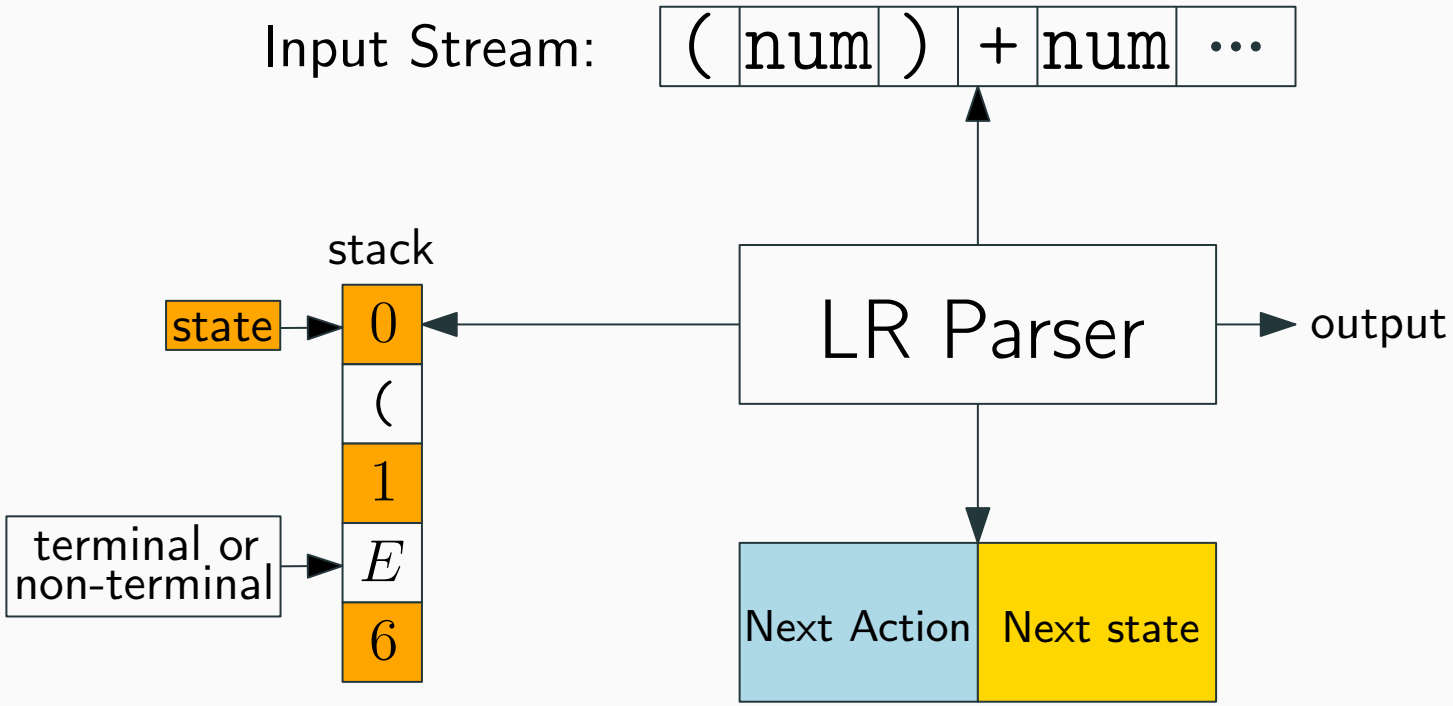
LR Parsing Table



Algorithm:

- Look at entry for current state S and input terminal C
- If Table[S,C] = s(S') then shift:
  - push(C), push(S')
- If Table[S,C] = $X \rightarrow \alpha$ then reduce:
  - pop($2 \times |\alpha|$), S'= top(), push(X), push(Table[S',X])

Model of LR Parser



LR(0) Parsing Tables

Algorithm for building LR(0) parsing tables:

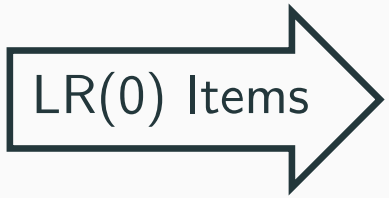
1. Compute parser states
2. Build a DFA to describe the transition between states
3. Use the DFA to build the parsing table

LR(0) States

- Each LR(0) state is a set of LR(0) items
- An LR(0) item is a production from the language with a separator somewhere in the RHS
- $X \to \alpha \bullet \beta$ says that
  - parser is looking for an X
  - it has an α on top of the stack
  - expects to find in the input a string derived from β

Example: LR(0) Items

r1	$S \to E\$$
r2	$E \to T$
r3	$E \to E + T$
r4	$T \to \texttt{num}$
r5	$T \to (E)$



1	$S \to \bullet E\$$
2	$S \to E \bullet \$$
3	$S \to E\$\bullet$
4	$E \to \bullet T$
5	$E \to T\bullet$
6	$E \to \bullet E + T$
7	$E \to E \bullet + T$
8	$E \to E + \bullet T$
9	$E \to E + T\bullet$
10	$T \to \bullet\texttt{num}$
11	$T \to \texttt{num}\bullet$
12	$T \to \bullet(E)$
13	$T \to (\bullet E)$
14	$T \to (E\bullet)$
15	$T \to (E)\bullet$

LR(0) Shift/Reduce Items

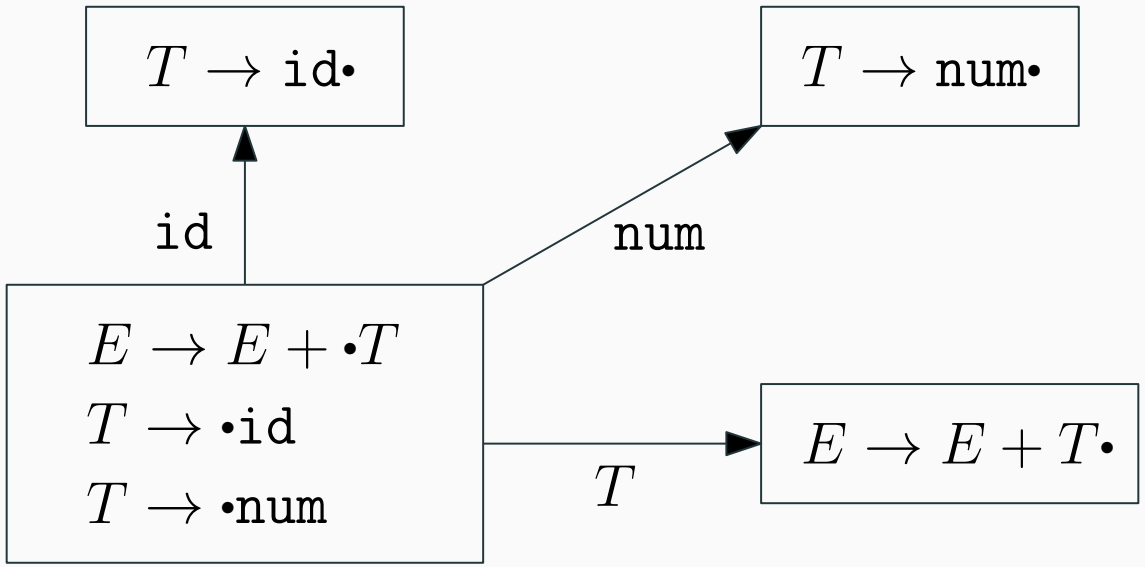
$N \to \alpha \bullet \beta$ Shift Item

$N \to \alpha\beta\bullet$ Reduce Item

LR(0) Automaton

An LR(0) automaton has:

- **states:** sets of LR(0) items
- **transitions:** label by grammar terminals or non-terminals



Construction of Automaton

- Augment grammar with production $S' \to S\$$
- Start state of automaton has empty stack $S' \to \bullet S\$$

To construct the automaton from the start state we need two functions:

- CLOSURE(L) to build its states
- GOTO(L, X) to determine its transitions

Begin with $\{S' \to \bullet S\$\}$, take the closure, and then keep applying GOTO

CLOSURE

If L is a set of items, CLOSURE(L) is the set of items such that:

- every item in L is in CLOSURE(L)
- if item $X \to \alpha \bullet Y \beta$ is in CLOSURE(L) and
  $Y \to \gamma$ is a production then
  $Y \to \bullet\gamma$
  is also in CLOSURE(L)

Exercise

- For the grammar

$$
\begin{aligned}
S &\to E \\
E &\to E + T \\
&\mid\ T \\
T &\to \texttt{num}
\end{aligned}
$$

  Compute the CLOSURE of the set of items $\{S \to \bullet E\}$

GOTO

If L is a set of items and X is a grammar symbol and $Y \to \alpha \bullet X \beta \in L$ then GOTO(L, X) is the CLOSURE of the set of all items $Y \to \alpha X \bullet \beta$

Exercise

- For the grammar

$$
\begin{aligned}
S &\to E \\
E &\to E + T \\
&\mid\ T \\
T &\to \texttt{num}
\end{aligned}
$$

  Compute the state that can be reached from the LR(0) state $\{S \to E \bullet\ ,\ E \to E \bullet + T\}$ on symbol $+$